

A Cost Improvement Rule Victimisation



J. Antony Vijay, Trishansh Srivastava, Manpreet Singh, Aman Kumar Soni

Abstract: Cloud computing is an effective technique used by developers and other users to implement different use styles. We will take up an issue in this paper which is very popular in the present era. “Storage-as-a-Service” is a category in cloud computing that uses cloud to store data with internet assistance. These online services use other categories of cloud computing namely “Infrastructure-as-a-Service, Platform-as-a-Service, and Software-as-a-Service”. “Storage-as-a-Service (StaaS)” is made available for the public use at certain prices. There are many service providers that are charging unreasonable amounts of money. So, we have come up with a better algorithm to overcome this problem. Our algorithm resolves the issue of transferring the data objects from one tier to another by excluding the prior information of the accesses. This algorithm allows us to find the breakthroughs for making transfer decisions. The proposed system makes sure that the Service Level Agreement (SLA) and reliability of reservations does remodel diverse resources and provide a tool that reduces victimization through Two-Tier StaaS.

Keywords: Cost Computing, StaaS, Cost driven optimization, Reduce, Reliability.

I. INTRODUCTION

Nowadays the standards of Storage-as-a-Service have gained thrust. According to IDC [1], there will be 175 ZB of volume of global data on the cloud. So, we need a more reliable and efficient way to store and access the data that is available on the cloud. The growing popularity of cloud services requires us to figure out more efficient and effective ways to store and share the data across different platforms at reasonable prices. There are many cloud storage providers that charge huge amounts of money or the prices are higher than they should charge. We can create an algorithm to tackle this problem. Through this algorithm we can actually make a difference in the cost that is required to access those services.

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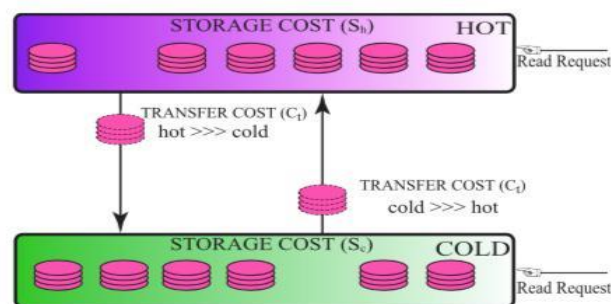
‘Storage-as-a-Service (StaaS)’ is made available, at certain rates, for public use. Most service providers charge excessive amounts of money. So, we came up with a better algorithm to solve this problem. Our algorithm solves the problem of moving the data objects from one tier to another by excluding the accesses prior information. This algorithm helps us to find the breakthroughs for making decisions about transfers.

Microsoft Azure, Google Cloud and Amazon AWS are the most preferred platforms that are used or trusted by the users all over around the globe. These service providers use multiple techniques to find out the cost that the users need to pay. There are many factors that affect the price of the cloud computing services. Storage size and access frequencies are the two of the major factors affecting the cost. We have created an algorithm to have an all-inclusive cost-driven optimization tool for using Storage as a Service (StaaS) clouds.

Cloud users can transmit the data and its objects to recover the price by fluctuating frequencies of access. Hence to make an appropriate transmission choice, they need to foresee the access graphs. However it is very challenging to foresee the graphs in case of cloud users

The current system uses an algorithm that enables us to pilot the “StaaS cloud” users to make decisions based on certain criteria. If there is any kind of cost reduction in transmission of data and the objects between “hot tiers” and “cold tiers” while there is no preceding information on coming access graphs.

The planned system makes a choice if it needs to transmit the data and its object in between the tiers using the internet without any information of the coming access graphs. The main step to make “an online cost optimization algorithm” is to manage to figure out the “break-even point” for making transmission choices.



II. LITERATURE SURVEY

A. “A Cloud-Oriented Content Delivery Network Paradigm: Modeling and Assessment”

This paper suggests, “Cloud-Oriented content delivery networks (CCDNs)” is made up of promising different traditional networks of content delivery.



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Disrupting the principles and the advantages of the cloud, like “pay as you go” “business model” and topographical distributing the sources, “CCDN” provides a wide and profitable result for networks of content transmission. In this paper, the “ranked framework” is discussed and concluded with an effective blend of distribution of content over a network of multi-provider cloud environments, where the overall cloud communication sources are together considered including the traditional cloud computing resources. “CCDN” replica placement issue within cloud also considers the limitations of the physical substrate.

B. “An Approach to Balance Maintenance Costs and Electricity Consumption in Cloud Data Centers”

This paper suggests that they deal with the issue of power consumption of the servers in Cloud Data Centre (CDC) and the cost for maintenance used to keep the servers’ processor running. They basically create a component consisting of many virtual machines and the requirements in the processor. Then they consider the prices involved in virtual machine processing on the servers, transmission of data objects within the virtual machines and the price for shifting the virtual machines across the servers. They have designed an algorithm to resolve this particular issue. The algorithm is “Maintenance and Electricity Costs Data Centre (MECDC)”

C. “Cost-Minimizing Dynamic Migration of Content Distribution Services into Hybrid Clouds”

This paper suggests that with the recent arrival of these technologies there is a growth in the distribution of content forcing the service providers for expandable and reasonable prices. The prime problem is to make sure that the services providers and their already existing applications provide better response time and value for money. Analysis depicts that the algorithm should have a response times that compliments the price that the users are paying the providers. They also say that these results are based on the prototype created by them.

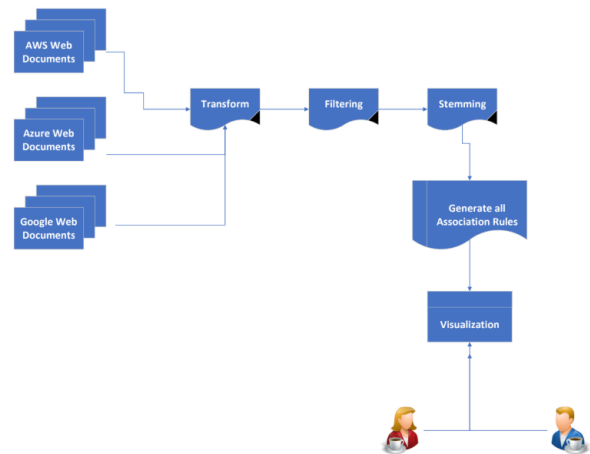
D. “Edge-CoCaCo: Toward Joint Optimization of Computation, Caching, and Communication on Edge Cloud”

This paper suggests that the enhancement of the innovative technologies in the form of different categories of realities and neural technology has extraordinary computing needs and graphic media is time sensitive. “Edge Cloud Computing (ECC)” has to be a legitimate solution to meet the demand of high activity. This paper a new path and the concept for computational solution. With caching policies. This algorithm is known as “Edge-CoCaCo ” which provides an optimum solution to the problem.

III. ARCHITECTURE

There are mainly three service providers that we have used to build the software using our algorithm. The service providers are Microsoft's Azure, Amazon's AWS and Google Cloud Services. We made an algorithm that considers the user/client requirement on the storage services provided by the above-mentioned service providers. We use certain criteria, a certain set of rules and regulations to narrow down the price range and which results in less victimization than there actually are. The software uses these

platforms and then transforms the data objects according to our needs. After this we pass the transformed data objects through a filter that narrows down the list of the data objects. Then the filtered data objects are streamed on to a platform that generates all the association rules. After going through all these steps, the output is then displayed on the users/clients’ window or the application. We call this step virtualization. Virtualization helps the user or the client to analyze the data objects in a much better and efficient way enabling the user or client to make better decisions whether to transfer the data objects from one tier to another without worrying about access protocols.



IV. MODULES

“Module 1: Amazon Data Extraction”

Sources show an “Object Oriented interface” to Amazon’s cloud services- AWS. They provide very good services in terms of abstraction and the calls by the clients. They have many methods and attributes that they have made available for the public use. They can further be categorized in actions, references, sub resources, etc. Boto3 only inherits the AWS APIs on the clients’ behalf. Boto3 provides 2 ways to access the above mentioned APIs:

1. “Client: low-level service access”
2. “Resource: higher-level object-oriented service access”

Conceptually, resources can be categorized into service resources (“sqs, s3, ec2, etc.”) and client resources (“sqs. Queue or s3.Bucket”). They neither have attributes nor identifiers. They share these components but don’t put them in use.

“An action is a method which makes a call to the service.” Actions puts the resources “identifiers” as their “parameters” executing a “low-level” resources but enables the user to add extra parameters using “keyword arguments”.

“Module 2: Azure Data Extraction”

“Azure SDK for Python” makes it easier for the users/clients to manage the codes from the application made on python. There are many libraries added to the “Azure SDK” to simplify the programming using Python. They now can upload any type of files, working with the tables and many other services that the programmer can make it for his/her use.

“Azure Command-line Interface (Azure CLI)” is a habitat to manage the libraries of Azure. Azure’s mainly focuses on the automation of the projects. The programmer needs to understand that there should be a hierarchical control line so as to protect the sensitive parts away from the public. Azure’s vision is to make an identity for the services and applications associated with Azure's resources. There should be a prohibited access area for the users in the applications to keep the confidential data secure. It suggests that there should be built in access for the users rather than creating a separate profile to access the sensitive data.

“Module 3: Cost Forecasts”

This cost model creates a link in between the source decision and the cost that the provider has spent on the development of the application and the maintenance of the servers. The cost for these services are set by keeping certain criteria in mind i.e. “resource cost, downtime during reconfiguration actions, and monetary Income associated with providing a particular level of service”. By using these techniques we are able to make an algorithm that considers “Service Level Agreement (SLA)”and maximum profits. There is a penalty when there is an allotment lower than the requirement. However, the penalty also depends on the “exponential penalty factor” and the “required value”.

V. DATASET DESCRIPTION

Long haul clients get the information in open SaaS mists are regularly secret. A large part of this information object is availed from 0-2 times every month and just a hardly any information objects are gotten to multiple occasions.

1) DATA DIVISION: To examine how our networked calculation performs under various outstanding tasks at hand designs, we order all information into 3 gatherings by their entrance change levels that are estimated as the proportion between σ and μ . We set the number of gets to an information object as being close to 5, so the standard deviation is additionally little. Assume the mean of an information object with a lifetime of one year is 0.167, so it is gotten multiple times in its lifetime.

2) COST: The earn back the original investment focuses β_c and β_h for each information object are determined at each time interim τ , in light of the fact that they are identified with the size of the information object and the quantity of aggregated read demands.

A. “ASSESSMENT OF OUR NETWORKED ALGORITHM”

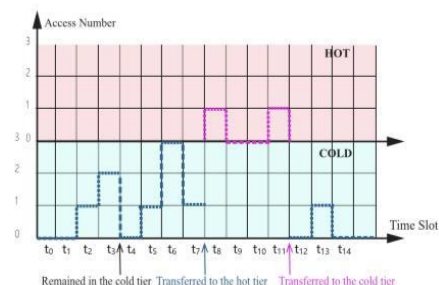
We presently assess the presentation of our networked calculation.

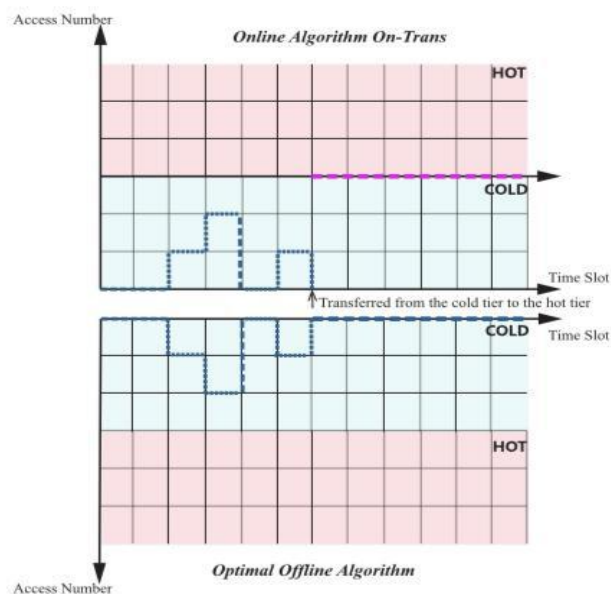
1) CRITERION ALGORITHMS

The main benchmark calculation is One-to-Trans, which takes various systems for information objects put away in the cold and hot levels. Calculation One-to-Trans moves information objects put away in the hot level to the chilly level when they have not been gotten to just because and moves information objects put away vulnerable level to the hot level once they get to. Furthermore, we utilize a third benchmark calculation, Off-Trans, which is a disconnected calculation and chooses the situation of information objects at each time interim τ with earlier information on future access recurrence in the next interim.

2) COST EFFICIENCY

Both networked calculations can accomplish cost reserve funds and the cost investment funds of calculation One-to-Trans are more than twice that of calculation One-to-Trans. The cost reserve funds of calculation One-to-Trans on class2 and class3 are not exactly zero, on the grounds that the information protests in class2 and class3 are perused rarely, they would be wise to be put away exposed to the harsh elements level for cost enhancement right now. Our networked calculation On-Trans has a superior exhibition on class2 and class3 since it can move information items to the hot level when they start to be gotten to every now and again, and the other way around. Our on the web calculation On-Trans has a superior outcome than calculation Stay-in-One and One-to-Trans whatever it runs on information objects which are gotten to every now and again or inconsistently. The legends On-Trans, Stay-in-One, One-to-Trans furthermore, Off-Trans individually speak to our networked calculation, Stay-in-One calculation, One-to-Trans networked calculation and Off-Trans disconnected calculation. When changing from Stay-in-One calculation to On-Trans calculation, beyond what 65% items can spare expenses and about 41% objects spare 10% of the expense. Furthermore, about 20% of the items can spare 20% of the expense by calculation On-Trans. One-to-Trans calculation likewise gets a decent presentation and about 32% items cut their expenses, yet just about 20% articles can spare 10% of the expense. When utilizing One-to-Trans calculation, there are just about 10% articles which cut their expenses, however over 85% items acquire a larger number of expenses than previously. When changing from Stay-in-One to our networked calculation On-Trans, about 60% items cut their expenses and about 59% articles spare over 2% cost. There are just about 35% objects reducing their expenses utilizing calculation One-to-Trans and over 60% items because more expenses. When changing from Stay-in-One to our proposed networked calculation On-Trans, about 70% of the items cut their expenses. The presentation of our networked calculation On-Trans is superior to anything calculation Stay-in-One. Contrasted and calculation Stay-in-One, about 22% of the articles spare over 10% of the expense, and about 10% of the items spare over 20% cost when clients pick our proposed networked calculation On-Trans. Practically all information items can reduce their expenses utilizing disconnected calculation Off-Trans. .





VI. RESULT

Consequently, whatever our online calculation On-Trans runs on information objects which are gotten to moderately often or inconsistently, or the amalgamation of these two sorts of information questions, its execution is better than online calculation One-to-Trans and calculation Stay-in-One.

VII. CONCLUSION

Right now, propose a networked calculation to direct SaaS cloud clients in settling on a choice on whether to move their information protests among cold and hot stockpiling levels for accomplishing cost improvements, while clients don't have to have any earlier information on future access frequencies to their information. In particular, right now use Google SaaS Cloud for instance to show our calculation. At that point, we demonstrate hypothetically that the proposed networked calculation can accomplish ensured serious proportions for information objects put away in a two-tier SaaS cloud. At last, through broad re-enactments we show that our networked calculation can accomplish critical expense reserve funds to SaaS cloud clients contrasted and continually keeping information protests in chilly levels or continually moving information objects to different levels when their entrance frequencies start to change. Right now, we have not considered composing activities that may change the size of information protests in our proposed on the web calculation. For future work, we will build up another on the web calculation which thinks about both peruse and composes solicitations to information questions in settling on moving choices while as yet ensuring serious proportions in sparring expenses.

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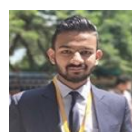
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