

Application of a Recurrence Matrix to Cryptography using Genetic Algorithm

Amit Kumar Mandle, Varsha Namdeo



Abstract: The aim of this paper is to develop an algorithm for encryption and decryption of a message involving recurrence matrix and genetic algorithm. In this approach secrecy is maintained by using of genetic algorithm. The proposed algorithm solves various problems that we are facing now days. Thus in this paper, with the help of recurrence matrix and genetic algorithm a secret sharing scheme is established, which is designed for encryption that maintains the secrecy of the information.

Keywords: Recurrence matrix, Genetic Algorithm, Crossover, Mutation, Encryption and Decryption.

I. INTRODUCTION

Due to increase of digital media transmission and unauthorized access of import data, secure data transmission over network has become a vital and critical issue. For information security cryptography uses mathematical techniques. The concept of encryption and decryption is the base of cryptography. When data is converted to some unreadable form, then it is called encryption, while that process is called decryption when unreadable form of data is again converted to its original form. Symmetric and asymmetric cryptography are two types of algorithm. In symmetric cryptography one key is used for encryption and decryption both, while in asymmetric cryptography there are two different keys are used, one key for encryption known as public key and other is decryption key known as private key [3].

A method which is based on natural selection and used for solving constrained and unconstrained optimization is known as genetic algorithm. There are three main types of rules that are used by genetic algorithm i.e. selection, crossover and mutation. In crossover technique combining two parents to form a children for next generation, while mutation process random changes the individual parents to form children. There are various types of crossover techniques i.e. single point, two-point, uniform etc. In this paper we will use two-point crossover technique. In this technique, we choose two random points on the chromosomes and genetic material

exchanged at these points. Mutation operation is also various types. Some mutation operations are insert, inversion, swap, flip, reversing, uniform etc. In this paper we will use swap mutation operation. In this operation we choose two bits random and swap their position.

When elements of a matrix are considered from a recurrence matrix then this type of matrix is known as recurrence matrix. In this paper we consider a recurrence matrix which is whose elements are taken from Fermat sequence 2,3,5,9,17,13,.....Thus in this paper we define a recurrence matrix made from Fermat Sequences follows:

$$R_F = \begin{bmatrix} 1 & C_{n+2} & C_{n+1} & C_n \\ C_{n+2} & 1 & C_{n+4} & C_{n+3} \\ C_{n+1} & C_{n+4} & 1 & C_n \\ C_n & C_{n+3} & C_{n+5} & 1 \end{bmatrix} = \begin{bmatrix} 1 & 5 & 3 & 2 \\ 5 & 1 & 17 & 9 \\ 3 & 17 & 1 & 33 \\ 2 & 9 & 33 & 2 \end{bmatrix}$$

II. LITERATURE REVIEW

Kumar [1], Surya [4], Patil [2] and others are developed various cryptographic techniques using algorithm along with other tools. They studied proposed easy cryptographic secure algorithm for communication. They also designed a secure communication method for encryption and decryption with the help of genetic algorithm.

III. METHODOLOGY

Following Kumar [1], Surya [4], Patil [2] and other researchers in this paper we develop and algorithm for encryption and decryption of a message using recurrence matrix and genetic algorithm we develop a secret sharing scheme for secure communication, which is designed in such a way that encryption maintains secrecy of the message.

IV. ALGORITHM

Numerical values for alphabets and some symbols used in the paper given in the following table:

Table – 1

A – 1	O – 15
B – 2	P – 16
C – 3	Q – 17
D – 4	R – 18
E – 5	S – 19
F – 6	T – 20
G – 7	U – 21
H – 8	V – 22
I – 9	W – 23
J – 10	X – 24

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K - 11	Y - 25
L - 12	Z - 26
M - 13	0 - 0
N - 14	

4.1 ENCRYPTION:

1. Convert the Plain text and arrange them in a block size of 16 bytes i.e. 4x4 order matrix.
2. Correct each alphabet of plaintext matrix into numeric value using Table I. Called this matrix M (say).
3. Multiply key matrix K(say) and Plaintext matrix M under modulo p, here we take p= 27, we get another matrix M1(say).
4. Convert the numeric value of each element of matrix M1 into 5-bit binary code and divided them into two segments.
5. Apply two point (agreed by sender and receiver) crossover technique.
6. Apply swap mutation operation on that positions which are shared with receiver.
7. Convert each decimal value of element of above resultant matrix into their corresponding alphabet/symbol using table I to get required cipher text.

4.2 DECRYPTION:

1. Consider the cipher text and convert each alphabet/symbol into corresponding numeric value using table I and arrange them into a matrix D (say) of order 4.
2. Convert numeric value of each element of matrix D into 5-bit binary form and divided them into two segments.
3. Apply swap mutation operation as shared with sender.
4. Applying two point crossover technique as shared by sender.
5. Convert each 5-bit binary group into their decimal equivalent and arrange them into a matrix of order 4. Say this matrix D1.
6. Multiply D1 and inverse of key matrix K under modulo p, here we take p = 27. We get another matrix say P.
7. Convert numeric value of each element of P into corresponding alphabet using Table I, to get plain text.

Illustration:

Encryption Steps:

1. Consider a recurrence matrix as key matrix of 4x4 order (non-singular), which is made from Fermat's sequence, as follows:

$$R_F = \begin{bmatrix} 1 & C_{n+2} & C_{n+1} & C_n \\ C_{n+2} & 1 & C_{n+4} & C_{n+3} \\ C_{n+1} & C_{n+4} & 1 & C_n \\ C_n & C_{n+3} & C_{n+5} & 1 \end{bmatrix} = \begin{bmatrix} 1 & 5 & 3 & 2 \\ 5 & 1 & 17 & 9 \\ 3 & 17 & 1 & 33 \\ 2 & 9 & 33 & 1 \end{bmatrix} \pmod{27}$$

$$= \begin{bmatrix} 1 & 5 & 3 & 2 \\ 5 & 1 & 17 & 9 \\ 3 & 17 & 1 & 6 \\ 2 & 9 & 6 & 1 \end{bmatrix} = K \text{ (say)}$$

2. Let a plain text be

Arrange them in a block size of 16 bytes i.e. 4x4 matrix, we get

$$\begin{bmatrix} P & UB & L \\ I & CE & N \\ C & RY & P \\ T & IO & N \end{bmatrix}$$

3. In the above block matrix substitute numeric values of letters using table I as follows:

$$M \text{ (say)} = \begin{bmatrix} 16 & 212 & 12 \\ 9 & 35 & 14 \\ 3 & 1825 & 16 \\ 20 & 915 & 14 \end{bmatrix}$$

4. Multiply the key matrix K and text Matrix M under modulo system, we get

$$M_1 = MK \pmod{27}$$

$$= \begin{bmatrix} 16 & 212 & 12 \\ 9 & 35 & 14 \\ 3 & 1825 & 16 \\ 20 & 915 & 14 \end{bmatrix} \begin{bmatrix} 1 & 5 & 3 & 2 \\ 5 & 1 & 17 & 9 \\ 3 & 17 & 1 & 6 \\ 2 & 9 & 6 & 1 \end{bmatrix} \pmod{27}$$

$$M_1 = \begin{bmatrix} 16 & 0 & 20 & 2 \\ 13 & 16 & 5 & 8 \\ 11 & 8 & 4 & 10 \\ 3 & 415 & 9 \end{bmatrix}$$

5. Convert the numeric value of each elements of matrix M1 into 5-bit binary code, we get

10000 00000 10100 00010 01101 10000 00101
01000
01011 01000 00100 01010 00011 00100 01111
01001

6. Apply two point crossover technique at 11 and 30, we get

10000 00000 00100 01010 00011 00100 00101
01000
01011 01000 10100 00010 01101 10000 01111
01001

7. Apply the swap mutation operation on 2nd and 4th position of bit in each 5-bit group, we get

10000 00000 00100 01010 01001 00100 00101
00010
01011 01000 10100 00010 01101 10000 01111
01001

8. Convert each 5-bit binary code into their decimal equivalent and arrange them in 4x4 matrix row wise, we get

$$\begin{bmatrix} 16 & 04 & 10 \\ 9 & 45 & 2 \\ 11 & 220 & 8 \\ 7 & 1615 & 3 \end{bmatrix} = C \text{ (say),}$$

9. Convert each decimal value into their corresponding alphabet/symbol using table I, we get the following cipher text

Decryption Steps:

1. Consider the cipher text
2. Convert each alphabet symbol into their corresponding numeric value using table I and arrange them into 4x4 square matrix D (say), we get -

$$D = \begin{bmatrix} 16 & 04 & 10 \\ 9 & 45 & 2 \\ 11 & 220 & 8 \\ 7 & 1615 & 3 \end{bmatrix}$$

3. Convert numeric value of each elements of matrix D into 5-bit binary form as follows:

10000 00000 00100 01010 01001 00100 00101
00010
01011 00010 10100 01000 00111 10000 01111
00011

4. Apply the swap mutation operation on 2nd and 4th position of bit in each 5-bit group, we get –
10000 00000 00100 01010 00011 00100 00101
01000

01011 01000 10100 00010 01101 10000 01111
01001

5. Apply two point crossover technique at 11th and 30th bit, we get –

10000 00000 10100 00010 01101 10000 00101
01000
01011 01000 00100 01010 00011 00100 01111
01001

Convert each 5-bit binary group into their decimal equivalent and arrange them into a matrix of order 4×4, we get –

$$\begin{bmatrix} 16 & 0 & 20 & 2 \\ 13 & 16 & 5 & 8 \\ 11 & 8 & 4 & 10 \\ 3 & 415 & 9 & \end{bmatrix} = D_1 \text{ (say)}$$

6. Now apply the operation

$$D_1 K^{-1} \pmod{27} = P \text{ (say)}$$

$$\Rightarrow P = \begin{bmatrix} 16 & 0 & 20 & 2 \\ 13 & 16 & 5 & 8 \\ 11 & 8 & 4 & 10 \\ 3 & 415 & 9 & \end{bmatrix} \begin{bmatrix} 0 & 1120 & 24 \\ 11 & 12 & 3 & 14 \\ 20 & 3 & 4 & 17 \\ 24 & 1417 & 22 \end{bmatrix} \pmod{27}$$

$$\Rightarrow P = \begin{bmatrix} 16 & 212 & 12 \\ 9 & 3 & 5 & 14 \\ 3 & 1825 & 16 \\ 20 & 9 & 15 & 14 \end{bmatrix}$$

7. Convert numeric values of each element of matrix P into corresponding alphabet using Table I, we get the final plain text as follows:

V. RESULT AND DISCUSSIONS

In this proposed algorithm, since we used recurrence matrix and genetic algorithm, therefore it is very difficult to break the cipher text without proper key. Extraction of original information from cipher text is very complicated because the chosen of recurrence matrix, secret key and genetic algorithm. Due to the size of key brute force attack is also hard task.

VI. CONCLUSION

Proposed algorithm is based on genetic algorithm and recurrence matrix. Here secrecy is maintained at following four levels:

1. Chosen recurrence matrix.
2. Chosen genetic algorithm.
3. Secret Key
4. Different Operations.

Therefore breaking of generated cipher text is a tedious job. Even when the algorithm is known, the extraction of plain text from cipher text becomes quite difficult.

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