

# Adjunct Octagonal Array Token Petri Nets



S. Kuberl, Anshu Murarka

**Abstract**— Adjunct Octagonal Array Token Petri Net Structures (AOATPN) are recently started out octagonal photograph delivering structures which out prolonged the Octagonal Array Token Petri internet systems. on this paper we view as AOATPN format over a control feature named inhibitor curves and separate it amongst a few critical octagonal photograph making and perceiving designs regarding the making energy.

**Index Terms**— Petri nets, octagonal array tokens, adjunction, octagonal grammars, octagonal tiling systems.

## I. INTRODUCTION

Hexagonal and hexagonal-recognized examples that appear to return to skip within the research and examination of the scene image cope with [12,13]. In [12] cluster hexagonal triangle society is seen as a three-dimensional depiction dimensional container, and "twin belief" of a given image collecting container. In biomedical picture makes sharpened, it's been examined that the device is programmed cells with hexagonal shape is commendable tool for short deal with snap shots of biomedical [9]. In a software software chromosome examination [9], the polygon surrounds associated with every picture appears as a hexagon. In view of the latter for the reason that Nineteen Seventies, formal mode to create or understand pictures of hexagonal laid out in writing [3,5-7,12,14] within the admission and examination machine of sample pics. a part of formalism antique fashion to create a cluster hexagonal Hexagonal Array SWIMMING Grammar (HKAG) [12] and hypothesis Hexagonal Array Grammars (HAG) [13]. Usefulness consecutive and brand new parallel development and catenations arrow point is the equal antique highlights of the fashion.

Hexagonal Tile Rewriting Grammars [16] and the nearby hexagonal Tile Rewriting Grammars [5] is stored without a doubt isometric tile based hexagonal shape version sentences, that have more regulations generative of HAG. Array hexagonal form Token net Petri (HPN) [7] has excelled at the string produces Petri nets [1, 4]. Petri internet is one of

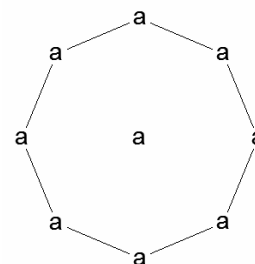
the traditional modes achieved to investigate the possible simultaneous framework, circulated, and parallel. In HPN, hexagonal clusters token is used to mimic the dynamics of the internet. In [7], the creators moreover provide a model of this hypothesis, Adjunct Hexagonal Array Token Petri net shape (AHPN), interest fuse adjunction, diffusion in state of affairs catenations sharp stones. model AHPN produce comparable institution of dialects made thru a part of the education HKAG and HAG. To accumulate the generative electricity, this time we want to keep in thoughts at some point of control highlight AHPN model, called inhibitors spherical section as in [8], the assessment and few flora and look at the picture of hexagonal expressive modes.

Pursuing the above experience we offer Octagonal Tile Rewriting Grammar and Language images [18]. This paper is taken care of out in a way that includes. In starting the phase, the importance of the octagon famous, nets Petri, and the thoughts of nets Petri overlay for cluster octagon terms chased that we study which means that from AOATPN and provide numerous fashions and then evaluation AOATPN and unique Octagonal display off the sentence shape ,

## II. COACHING

*Definition 2.1:*

permit  $\square$  be meeting the letter is restrained in an effort to snap shots. A photo of the side  $\square$  p octagon is an octagonal picture cluster of  $\square$ zero. Octagons indicated inside the accompanying species:



**Figure 2.1**

For example, an octagonal picture over the alphabet  $\{a\}$  is

*Definition 2.2:*

The collection of octagonal arrays along with the alphabet  $\Sigma$  is denoted by  $\Sigma_0^{**}$ . An octagonal picture language L over  $\Sigma$  is a subcollection of  $\Sigma_0^{**}$ .

*Definition 2.3:*

For  $p \in \Sigma_0^{**}$ , allow  $\hat{p}$  is the octagonal array hold by surrounding p with a special boundary symbol  $\# \notin \Sigma$ , for example

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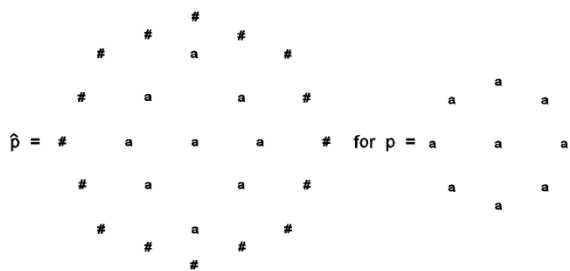


Figure 2.2

Definition 2.4:

The four tuple (a, b, c, d) is denotes the length of the picture p denoted by  $|p| = (a, b, c, d)$ . Let  $p_{ijk\ell}$  be the symbol in p with co-ordinates (i, j, k,  $\ell$ ) where  $1 \leq i \leq a, 1 \leq j \leq b, 1 \leq k \leq c, 1 \leq \ell \leq d$ . Let  $\Sigma_0^{(a,b,c,d)}$  be the collection of octagonal pictures of length (a, b, c, d).

Definition 2.5:

Given an octagonal picture p of length (a, b, c, d), we called by  $B_{e,f,g,h}(p)$  the collection of all octagonal subpictures of p of length (e, f, g, h), where  $e \leq a, f \leq b, g \leq c, h \leq d$ . Every portion of  $B_{2,2,2,2}(p)$  is called an octagonal tile.

Definition 2.6:

A non-convex octagon ABCDEFGH as shown in Fig. 1 is called an arrowhead if  $|BC| = |HA|, |CD| = |GH|, |DE| = |FG|$ , BC is parallel to AH and CD is parallel to GH and DE is parallel to FG. It is noted that the opposite sides  $|AB|$  and  $|EF|$  are equal and parallel.  $|AB|$  is the thickness of the arrowhead. BCDE is the outermost edge and FGHA is the innermost edge.

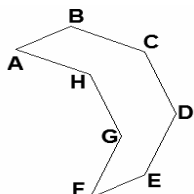


figure 2.3 proper better pointed stone

relying at the manner of a vertex a pointed stone is surveyed as higher sharpened stone (U), proper top sharpened stone (RU), right pointed stone (R), right diminishing pointed stone (RL'), decline sharpened stone (L'), left lower pointed stone (LL'), left sharpened stone (L), left top sharpened stone (LU).

Definition 2.7:

within the event that PQRSTUWV (Fig. 2.) is an octagon and ABCDEFGH is a pinnacle valid sharpened stone, at that aspect the sharpened stone may be catenated to the octagon interior the appropriate top heading.

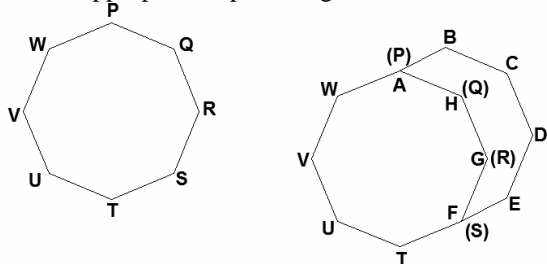


Figure 2.4

It is noted that  $|WB| = |WA| + |AB| = |WP| + |PB|$  and  $|TE| = |TF| + |FE| = |TS| + |SE|$ .

Similarly the remaining arrowhead catenations can be defined.

### III. OCTAGONAL ARRAY TOKEN PETRI NETS(OATPN) & RESULTS

Definition 3.1:

A petri net structure is a four tuple  $C = (P, T, I, O)$  where  $P = \{p_1, p_2, \dots, p_n\}$  is a finite collection of places,  $n \geq 0$ ,  $T = \{t_1, t_2, \dots, t_m\}$  is a finite collection of transitions  $m \geq 0$ ,  $P \cap T = \emptyset$ ,  $I: T \rightarrow P^\infty$  is insituate function from transitions to bags of places and  $O: T \rightarrow P^\infty$  is Outsituate feature of the transition into the wallet of vicinity.

Definition 3.2:

A Petri net marking is a mission of tokens to locations petri nets. Tokens are required to decide the implementation of Petri nets. the quantity and feature of the token may be changed throughout the execution petrinet a.

Definition 3.3:

Inhibitor arc from  $p_i$  to transition  $t_j$  has a small circle on the arrow within the bow normal. This approach that the transition  $t_j$  is enabled first-class if  $p_i$  does not have a token. Transition is activated handiest if all the ordinary insituates tokens and all insituates inhibitor that has zero tokens.

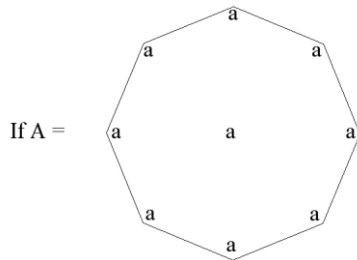
in this paper an octagonal array over the alphabet used as a token.

Definition 3.4:

If  $C = (P, T, I, O)$  is a Petri internet structure with octagonal array more than the mark as the beginning ,, label at the least one transition into catenation guidelines arrows and a limited series of prevent factor, then a Petri net C shape is defined as Octagonal Array Token shape Petri net (OATPNS).

Definition 3.5:

If C is OATPNS the language generated by using way of Petri internet C is described as  $L(C) = O / O$  in P for a few P in F with the array in a specific area as a prelude that indicates each single practicable succession exchange is permanent. collecting each single octagonal exhibition in very last location F is called a language created by using C.



the number of columns of A is 2, firing t adds upper arrowhead catenation A<sub>1</sub>, we get B

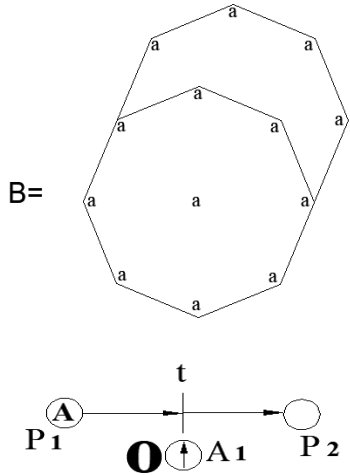


Figure 3.1 Position of token before firing

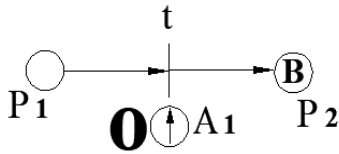
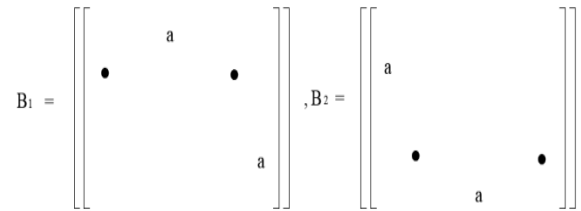
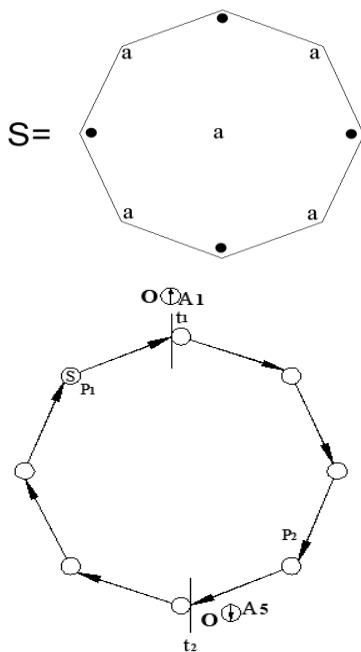
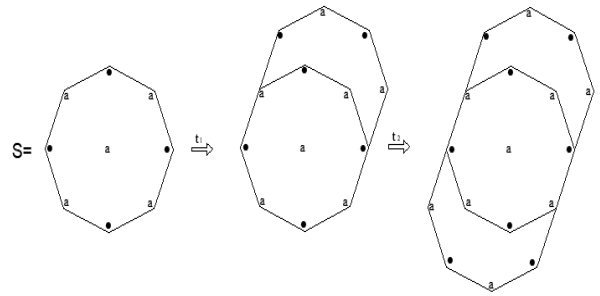


Figure 3.2 Position of token after firing

Example 3.1  $\Sigma = \{a, \bullet\}$ ,  $F = \{P_1\}$



Terminating t<sub>1</sub> arranges an exhibit in P<sub>2</sub> making t<sub>2</sub> enabled. Firing t<sub>2</sub> situates a cluster in P<sub>1</sub>. The terminating succession (t<sub>1</sub>t<sub>2</sub>)<sup>k</sup>, k ≥ 0. At the point when the changes t<sub>1</sub>, t<sub>2</sub> fire the exhibit that compasses the outside spot is appear below.



The language generated by this OATPNS is Octagonal Rangoli.

#### IV. ADJUNCT OCTAGONAL ARRAY TOKEN PETRI NET STRUCTURE (AOATPN)

Adjunction is a induction of arrowhead catenation. Using the upper arrowhead (U) catenation  $\textcircled{\uparrow} X$ , the arrowhead X is catenated to O after the unit upper arrowhead present in the boundary of O. However an upper adjunction be able to join the array X into array O after or before any unit upper arrowhead of O.

Let O be an octagonal array of length  $(|O|_a, |O|_b, |O|_c, |O|_d)$  in  $X^*$  called anchor array;  $X \in X^*$  be an arrowhead language whose members, named adjunct arrow heads, have permanent thickness and changeable length which depend on the consequent length of parameters of the anchor array O. For example, if X is an adjunct upper arrowhead,  $|O|_a$  (thickness) is permanent and the other three parameters  $|O|_b, |O|_c, |O|_d$  (length) depend on the consequent parameters  $|O|_b, |O|_c, |O|_d$  of the anchor array O.

In an anchor array O, there are  $|O|_a$  number of unit upper arrowheads (U) (lower arrowheads (L')) at hand, which we denote by  $u_1, u_2, u_3, \dots, u_{|O|_a}$  ( $l'_1, l'_2, l'_3, \dots, l'_{|O|_a}$ ). Here,  $u_1(l'_1)$  denotes the boundary unit arrowhead and  $u_{|O|_a}(l'_{|O|_a})$  denotes the innermost unit arrowhead in the U (L') direction. Any position between  $u_i$  and  $u_j$ ,  $i < j$ , is called after  $u_i$  ( $au_i$ ) or before  $u_j$  ( $bu_j$ ). An U (L') adjunct arrow head X can able to be connected into the anchor array O in  $|O|_a + 1$  positions subject to the condition of arrowhead catenation. An U (L') adjunction rule is a tuple  $(O, X, bu_i(b'l'_i) / au_i(al'_i))$ ,  $1 \leq i \leq |O|_a$  connecting X into O before  $u_i(l'_i)$  or after  $u_i(l'_i)$ .

Similarly, in a anchor array O,  $|O|_b$  a number of unit arrowheads in the right upper(RU) (left lower(LL)) direction are found. They are denoted by  $ru_1, ru_2, ru_3, \dots ru_{|O|_b}(ll_1, ll_2, ll_3, \dots ll_{|O|_b})$ . An RU (LL) adjunct arrowhead X can able to be connected into the anchor array O in  $|O|_b + 1$  positions subject to the condition of arrowhead catenation. An RU (LL) adjunction rule is a tuple  $(O, X, bru_i(bll_i) / aru_i(all_i))$ ,  $1 \leq i \leq |O|_b$  connecting X into O before  $ru_i(ll_i)$  or after  $ru_i(ll_i)$ .

In the anchor array O, there are  $|O|_c$  number of unit right arrowheads(R) (left arrowheads(L)) present, which we denote by  $r_1, r_2, r_3, \dots r_{|O|_c}(l_1, l_2, l_3, \dots l_{|O|_c})$ . An R(L) adjunct arrow head X can able to be connected into the anchor array O in  $|O|_c + 1$  positions subject to the condition of arrowhead catenation. An R(L) adjunction rule is a tuple  $(O, X, br_i(bl_i) / ar_i(al_i))$ ,  $1 \leq i \leq |O|_c$  connecting X into O before  $r_i(l_i)$  or after  $r_i(l_i)$ .

Again, in a anchor array O,  $|O|_d$  a number of unit arrowheads in the right lower(RL) (left upper(LU)) direction are found. They are denoted by  $rl_1, rl_2, rl_3, \dots rl_{|O|_d}(lu_1, lu_2, lu_3, \dots lu_{|O|_d})$ . An RL(LU) adjunct arrowhead X can able to be connected into the anchor array O in  $|O|_d + 1$  positions subject to the condition of arrowhead catenation. An RL(LU) adjunction rule is a tuple  $(O, X, brl_i(blu_i) / arl_i(alu_i))$ ,  $1 \leq i \leq |O|_d$  connecting X into O before  $rl_i(lu_i)$  or after  $rl_i(lu_i)$ .

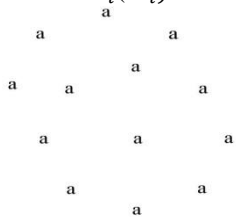


Figure 4.1

Figure 2.1 shows all the unit arrowheads in the U,RU,R,RL directions for the octagon in Figure 4.1.

Definition 4.1:

An Adjunct Octagonal Array Token Petri Net Structure(AOATPN) is a five tuple  $Q = \{\Sigma_0, C, M_0, \rho, F\}$  where  $\Sigma_0$  is a given alphabet,  $C = (P, T_0, I, O)$  is a petri net structure [7,8,19] with tokens as octagonal arrays over  $\Sigma_0$  and  $T_0$  contains transitions with inhibitor arcs,  $M_0: P \rightarrow \Sigma_0^{**}$ , is the preliminary marking of the net  $\rho: T_0 \rightarrow L$ , A mapping of opportunity series for stacking name that some development also can let alone have validated adjunction stone regulations for names and  $F: P$ , is the buildup of confined spots conclusive. In AOATPN, style changes that could empower and ends like that of Octagonal Petri net [19] aside from the type (iii) wherein the names of progress pointy rocks adjunction control in lieu of tips catenation sharpened stones.

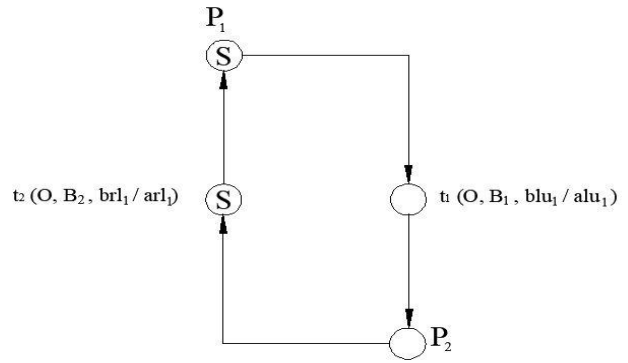
Definition 4.2:

in the occasion that Q is AOATPN, the trouble that the language of snap shots octagonal made thru Q inferred as in location for numerous q q in F. commencing with an octagonal array (token) is determined because the initial sign of the alphabet, all sequences transition opportunity fired. the following series of all arrays in places surrender F named language created by means of AOATPN. We have been

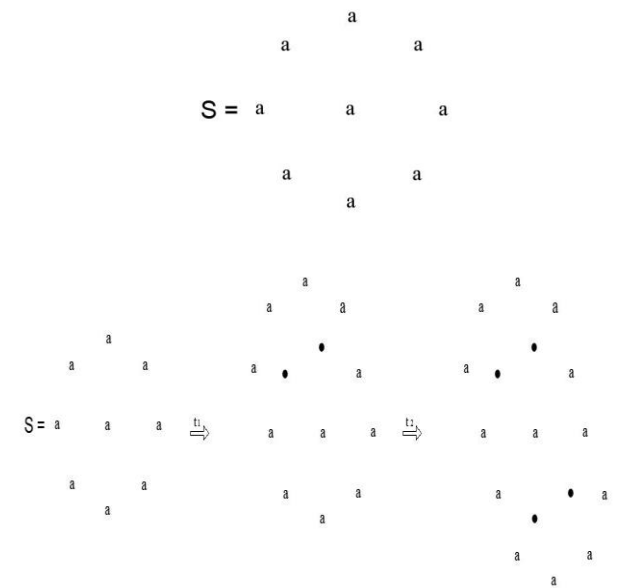
tested by manner of the language institution AOATPNL octagon photograph produced via using Adjunct Octagonal Array shape Token Petri internet.'

Example 4.1:

Consider the AOATPN,  $Q_1 = \{\Sigma_0, C, M_0, \rho, F\}$  where  $\Sigma_0 = \{a, b\}$ ,  $C = (P, T_0, I, O)$ ,  $P = \{p_1, p_2\}$ ,  $T_0 = \{t_1, t_2\}$ ,  $I(t_1) = \{p_1\}$ ,  $I(t_2) = \{p_1\}$ ,  $M_0$  is the preliminary marking; the array S is in  $p_1$  and there is no array in  $p_2$ .  $\rho(t_1) = (O, B_1, blu_1)$ ,  $\rho(t_2) = (O, B_2, brl_1)$  and  $F = \{p_1\}$ .



The arrays are used in the way of following:



At starting  $t_1$  is the only one enabled transition. Firing of  $t_1$

adjoins LU arrowhead 'a' inside the limit LU pointed stone of S and arranges the new cluster O in the spot of q2, making  $t_2$  empowered. Terminating  $t_2$  appends a RL

pointed stone 'a' inside the limit RL pointed stone of the exhibit O in  $p_2$  and arranges the new and last cluster in  $q_1$ . Refresh this equivalent terminating rules upto n times we will get an octagonal cluster pursues:





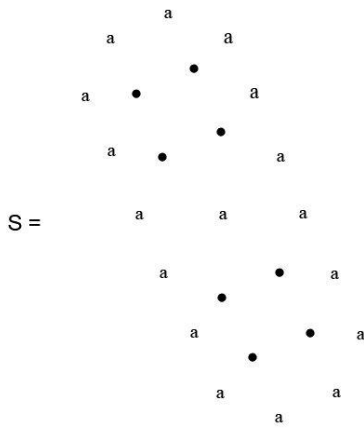
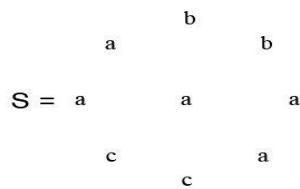


Figure 4.2 Octagonal Array in  $L_1$

**Example 4.2:**

The AOATPN  $Q = \{\Sigma_0, C, M_0, \rho, F\}$  with  $\Sigma_0 = \{a, b, c\}$ ,  $F = \{f_1, f_2\}$  given in the figure 5.3. Where



$\langle \begin{smallmatrix} a & a \\ c & b \end{smallmatrix} \rangle A \langle \begin{smallmatrix} b \\ a \end{smallmatrix} \rangle$ , generates the language  $L_2$  of an octagonal

arrays of length  $(2m, 2, 2, 2)$ ,  $m \geq 1$ , with inner elements over with the  $x_2$  direction and opposite of  $x_2$  direction forming the pattern  $b^{2m}c^{2m}$  and the remaining (direction) boundary elements are 'a's.

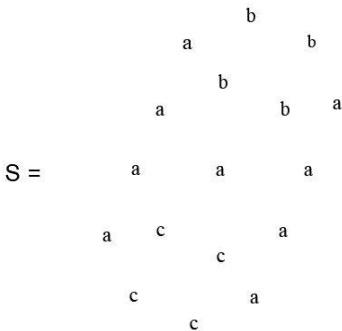


Figure 4.3 Octagonal Array in  $L_2$

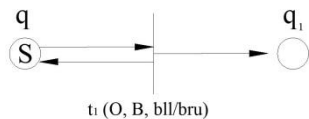


Figure 4.4 Petri net to generate  $L_2$

**Definition 4.3:**

A pure 2D octagonal context-free grammar ( $P_{2D}OCFG$ ) [14,15] is  $G = \{\Sigma_0, P_U, P_L, P_{RU}, P_{LL}, P_L, P_R, P_{RL}, P_{LU}, M_0\}$  where  $\Sigma_0$  is a finite collection of symbols,  $P_u = \{t_{ui} / 1 \leq i \leq m\}$ ; each  $t_{ui}$  is called a U table, is a collection of context-free rules of the form  $b \rightarrow \beta$ ,  $b \in \Sigma_0, \beta \in \Sigma_0^{**}$  such that any two rules of the variety  $b \rightarrow \beta, c \rightarrow \gamma$  in  $t_{ui}$ , we get  $|\beta| = |\gamma|$ , there  $|\beta|$

denotes the length of  $\beta$ . Similarly define the remaining seven components  $P_L, P_{RU}, P_{LL}, P_L, P_R, P_{RL}, P_{LU}$ .  $M_0$  is a finite collection of axiom array that are octagonal arrays.

Derivations are defined as follows:

For any two octagonal arrays  $O_1, O_2$ , we write  $O_1 \Rightarrow O_2$  if  $O_2$  is obtained from  $O_1$  by rewriting all the symbols in an unit arrowhead of  $O_1$  by rules of a relevant table in  $P_U \cup P_L \cup P_{RU} \cup P_{LL} \cup P_L \cup P_R \cup P_{RL} \cup P_{LU}$ .

$\Rightarrow^*$  is the reflective transitive closure of  $\Rightarrow$ .

The octagonal picture language  $L(G)$  generated by  $G$  is the collection of  $\{O/O_0 \Rightarrow^* O \in \Sigma_0^{**}, \text{ for some } O_0 \in M_0\}$ .

The group of all octagonal picture array languages generated by pure 2D octagonal context-free grammars is denoted by  $P_{2D}OCFL$ .

**Definition 4.4:**

A pure 2D octagonal context-free grammar with regular control ( $P_{2D}^R OCFG$ ) is a tuple  $Gr = (G, \Gamma, C)$ , where

- (1)  $G$  is  $P_{2D}OCFG$ ,
- (2)  $\Gamma$  is the control alphabet, the collection of labels of the rule tables in  $P_U \cup P_L \cup P_{RU} \cup P_{LL} \cup P_L \cup P_R \cup P_{RL} \cup P_{LU}$ ,
- (3)  $C \subseteq \Gamma^*$  is the regular control associated with the  $Gr$ .

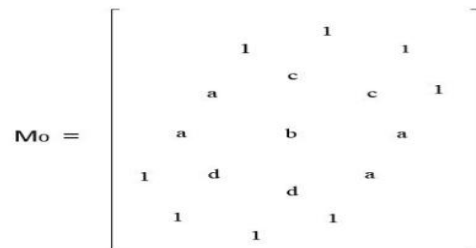
If  $O \in \Sigma_0^{**}$  and  $O_0 \in M_0$ ,  $O$  is derived from  $O_0$  in  $Gr$  by means of a control word  $v = v_1 v_2 \dots \in C$ , in symbols  $C \Rightarrow_v O$ , if  $O$  is obtained from  $O_0$  by applying the table rules as in the sequence of tables  $v = v_1 v_2 \dots$ . The language  $L(G)$  generated by  $P_{2D}^R OCFG$ .  $Gr$  is the collection of pictures  $\{O/O_0 \Rightarrow_v O \in \Sigma_0^{**}, \text{ for some } v \in C\}$ . The collection of all octagonal context-free grammars with regular control is denoted by  $P_{2D}^R OCFL$ .

**Definition 4.5:**

A 2d natural octagon context-free grammar spacious  $Exp_{2D}OCFG$  normal controls are second octagonal context-unfastened grammar herbal with everyday controls, collectively with the alphabet that is alphabetical terminate a good sized picture of an photograph octagon and Has collected manage pictures involved sufficient in induction strategies and they may now not seem inside the very last image. the accumulation of all dialects octagonal pix created thru growing natural placing octagonal unfastened penalty tool with ultra-current second manage validated through the use of  $Exp_{2D}OCFL$ .

**4.3 version:**

The language  $L_2$  in version 4.2 may be produced through  $Exp_{2D}OCFG (G, \{ru_1, ru_2\}, C)$  where  $G = (\mathcal{T}^f \cup \mathcal{T}^c, \{ru_1, ru_2\}, \{v_n, n \leq 8\}, M_0)$  and  $\mathcal{T}^f = \{1, 2, 3, \bullet\}$  and  $\mathcal{T}^c = \{a, b, c, d\}$



$ru_1 = \{a \rightarrow 1a, d \rightarrow 2d, d \rightarrow 2d, a \rightarrow 1a\}$   
 $ru_2 = \{a \rightarrow 1a, c \rightarrow 3c, c \rightarrow 3c, a \rightarrow 1a\}$  and the control language  $C = \{((ru_1)(ru_2))^*((ru_1')(ru_2'))\}$  where  $ru_1' = \{a \rightarrow 1, d \rightarrow 2\}$  and  $ru_2' = \{a \rightarrow 1a, c \rightarrow 3, c \rightarrow 3c, b \rightarrow 1\}$

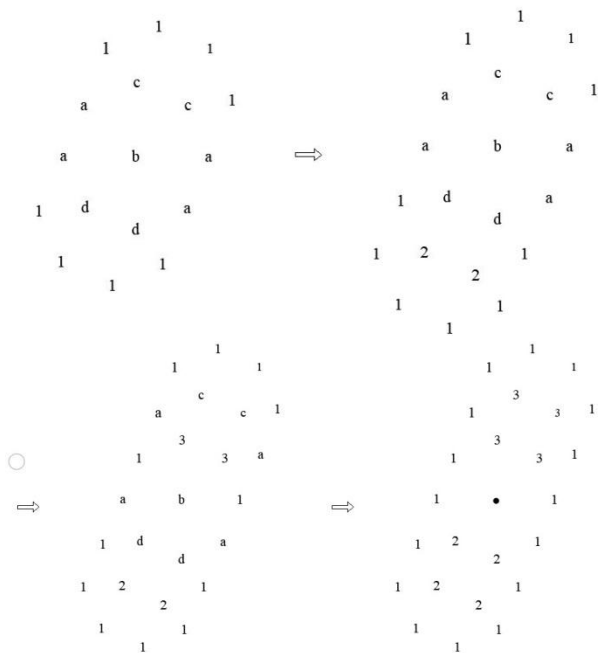


Figure 4.4 Octagonal Array in  $L_3$

*Theorem 4.1:*

The proper circle of relatives PR2DOCFL contained in Octagonal Adjunct Array Token Petri internet Languages (AOATPL).

affirmation:

take into account L into language octagon picture created through a in easy terms second-octagonal arrangement freed from punctuation  $G = \square$  zero, PU, PL', PRU, PLL, PL, PR, PRL, PLU, extra languages joint customs control amassing names (RU1, RU2, ..., rum).

in the table RU, like Trui, patch up rui gabled stone rui then all the snap shots to be changed thru a way for attention of fiction in parallel. Regardless, inside the concept of the RU-adjunction rule (O, B1, Arui / brui) will in addition have a very last product this is just like that of a large Trui table. In what's the foundation inside the back of the time period is offensive RU RU desk adjunction.

along those traces, for all the paintings region Trui RU, RU core adjunction relative fee can be described.

4.2 idea:

Exp2DOCFL and AOATPL now not disjoint.

affirmation:

AOATPL language L2 do not forget the four.2 fashions introduced are in Exp2DOCFL (assessment and 4.three shape). alongside pressure Exp2DOCFL and AOATPL not disjoint.

prevent

Paper ponder a change gathering Octagonal shape Array Token Petri nets more adjunction sprucing stone because the call implies alternate, close to spotlights bend controls known as inhibitors. We associated this version with exceptional style octagonal shape PR2DOCFG sentence, and

Exp2DOCFG. we have found out that AOATPN have step by step better confinement of PR2DOCFG but unmatched and non-decipher with a totally massive designs. there may be not anything in any respect in anyway between AOATPN and OATPN complexity.

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