

Design and Implement a Real-Time Detection and Defence Mechanism Against the SYN Flood Attack in Server Client System



SumontoSarker, KritimoyBosu, Firdous Bin Ismail, Md. MahabubHossain, Md. Mehedi Islam

Abstract: In the digital world, maintaining information is much difficult. Without security measures and controls in place, data might be subjected to an attack. Now a day's several attacks are evolved & Distributed Denial of Service (DDOS) is one of them. There are various categories of DDOS attack. SYN flood is addressed as one of the most dangerous attacks. In three way handshaking method a SYN packet is generated and a received ACK acknowledgement is provided to the corresponding. When the SYN packet is generated continuously from random sources is considered as flooding. And it's known as SYN flood attacks. This paper is constructed with a proposed technique for the betterment of both the detection and defense techniques against it. The detection process is improved by a database added in the server for accepting random flooding for a limited time interval. And the defense algorithm is a developed design operated by scrolling the pending requests from database and checking the accessibility of the user and stop requesting otherwise. There are two parts of this research paper. The first one is to discuss extensively the various aspects of SYN Flood attack and developing the knowledge of this flooding attack mechanisms and the second one is to detect the SYN Flood attack and finding a better mitigation process through which we can reduce the loss of any information that generally happens by this deadly flooding. For saving server from crush it is important to mitigate this attack. So it may prove effective in home appliance servers like IoT, IoE that any of the fraud can't get access into the server for any harmful activity.

Keywords: Botnet, mitigation, TCP SYN flooding, server client system attack, defense method.

I. INTRODUCTION

Nowadays, distributed Denial of Service (DDoS) attacks pose one of the most serious security threats to the Internet. DDoS attacks can result in a great damage to the network service according to [1][2]. Network and data are vulnerable to network attacks which may include DDoS attacks launched by attackers around the world to disrupt the network. DDoS attacks are categorized as the most popular network attack because the attacks are most common around the world. [3][4]. As DDoS attack is easy to implement and its attack method is simple, so it's spreading incrementally day by day. But yet it's difficult to defend this problem. The DDoS attackers usually utilize a large number of puppet machines such as construct a BOTNET to launch attacks against one or more targets, which can exhaust the resources of the victim side that makes the victim lose the capability to serve legitimate customers and prevent legitimate users from accessing information or services as shown in the figure 1. Since DDoS attacks can greatly degrade the performance of the network and are difficult to detect, they have become one of the most serious security challenges to the current intrusion detection systems (IDS). However, as long as the flooding is detected early, the loss can be reduced to the minimum. Therefore, DDoS attack detection and defence still attract much concern from researchers. The research is on the detection of SYN Flood attack and the prevention of it. For these types of detection and defence mechanism the system maintains the scenario of being filtered to all the requests to remain the server free from unusual flooding as shown in the Figure 1 with the broken arrows. The SYN Flood attack related previous studies, implementation of it in a single or multiple client, detection of it and the defence against it described with the association of the flowchart and recovery process is also attached with associated figures and then the finding out put simulation and the results are also figured graphically. Entire process of simulation can be clarified with the following diagrammatic process.

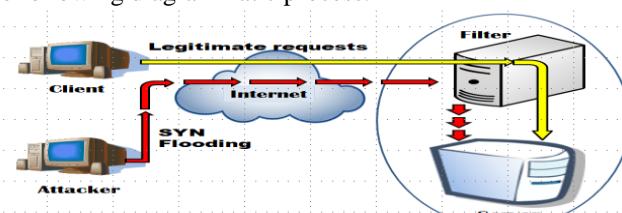


Fig. 1. SYN Flood Detection and Prevention Model

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* Correspondence Author

SumontoSarker*, Department of Electronics and Communication Engineering of Hajee Mohammad Danesh Science and Technology University, Dinajpur, Bangladesh. Email: sumonto@hstu.ac.bd

KritimoyBosu, Department of Electronics and Communication Engineering of Hajee Mohammad Danesh Science and Technology University, Dinajpur, Bangladesh. Email: bosukritimoy@gmail.com

Firdous Bin Ismail, Department of Electronics and Communication Engineering of Hajee Mohammad Danesh Science and Technology University, Dinajpur, Bangladesh. E-mail: bishalhstu@gmail.com

Md. MahabubHossain, Department of Electronics and Communication Engineering of Hajee Mohammad Danesh Science and Technology University, Dinajpur, Bangladesh. Email: Im.mahabub@gmail.com

Md. Mehedi Islam, Department of Electronics and Communication Engineering of Hajee Mohammad Danesh Science and Technology University, Dinajpur, Bangladesh. E-mail: mehedi@hstu.ac.bd

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Formatting an active mechanism its much important to know about the TCP (Transmission Control Protocol)[5]. It is a collection of rules, it and the Internet Protocol (IP) is used together to exchange information between computers via the internet send data in the form of a unit.

The IP protocol controls the actual data transmission, the TCP protocol is primarily responsible for tracking the distribution of information transmitted over the internet, individual data units (packages).[26] Establishing TCP requires three handshaking to build first, the client sends a connection request message, and the Server segment accepts the connection. It then replies with an ACK message and allocates resources for this connection. Client after receiving the ACK packet, the terminal also generates an ACK packet to the Server segment. And allocate resources so that a TCP connection is established. [6] At the server's end most of the time security threats are the major facts. SYN flooding is one of these.

A. SYN Flood attack

SYN Flood is currently the most popular DoS (Denial of Service Attack) and One of the ways of DDoS (Distributed Denial of Service Attack), this is a send a large number of forged TCP connection requests using TCP protocol flaws[23], Commonly used fake IP or IP number segments sent the first request for a massive number of requests. The handshaking packet (SYN packet) when requests the server responds to the second handshake packet by (SYN+ACK package). If the other party has a fake IP, then client will never receive the package and server will not respond to the third handshake process [5]. Clients cause the server a large number of SYN_RECV and retry the default 5 times for the second handshake packet, stuffed with TCP waiting for the connection queue [25]. When Source is exhausted (CPU full or out of memory) for normal business requests the real connection does not come in. Due to the principle of SYN flood attack, the corresponding monitoring is also very simple [20]. There are two popular methods on the network, namely: The first is to shorten the SYN Timeout time due to the SYN Flood attack and the effect of hit depends on the number of SYN semi-joins maintained on the server [27]. Value = the frequency of the SYN attack x SYN Timeout, so by shortening from Receiving a SYN packet to determine that the packet is invalid and discard the connection. Time, for example, set to 20 seconds or less, can double the server the load, however, too low SYN Timeout settings may affect customers [6].Normal access the second method is to set the SYN cookie,just give each one please Assign a cookie to the connected IP address, if it is continuous for a short time a repeated SYN message of an IP is considered to be an attack. And record the address information, the package from this IP address will be thrown away. The result of this problem may also affect the access of normal users [19]. According to [13] the above two methods can only deal with the original SYN Flood. Attack; shorten SYN Timeout time only when the attack frequency of the other party is not high In the case of the SYN Cookie, the SYN Cookie is more dependent on the other party's use of the real IP address, if an attacker sends a SYN message at hundreds or thousands of requests/second. At the same time, it uses SOCK_RAW to randomly rewrite the source address in the IP packet at servers end.

II. PREVIOUS APPROACHES OF DETECTION

Information may have many resources to be grabbed. But the most used and reliable source of information. in the modern world is internet. It's the blessing of internet that it serves news feed and many more across the world. Besides it needs to be protected and also information to be secured. Security of information has three components these are confidentiality to protect against unauthorized individuals, integrity for protection against alteration, and availability for the protection against the interference with the means to access the resources as described in [7] Many organization implant the firewall protection barrier around the internet. Networks are much complicated to defend using only traditional techniques such as cryptographic techniques stenographic technique authentication and static firewall. Many methods have proposed for the detection of this SYN flooding attacks. The available techniques are able to find out the attack. One of them is hash table. It's a high-performance method to detect attack that's based on IP trackback. Generally the DoS attacks can be traced through three classified processes [14]. These are (1) Based on router Data Structure,(2) Statistical analysis of packet flow and (3) Fuzzy logic and neural network known as AI (Artificial Intelligence).

A. Router Based Detection Scheme

Bloom filter is a space efficient data structure used in a router for pattern matching in many network communications. It is used to inspect packets based on many algorithms[7]. The main advantage of this technique is it uses the change point detection method based on nonparametric cumulative sum (CUSUM) for the retransmission of SYN packets [8]. It also presents a cumulative analysis result based on Counting Bloom Filter (CBF) in leaf router. The CBF is used to store the full information for TCP connection which includes client server IP address ports and initial sequence numbers [8]. But it has disadvantages also as its inefficient if in case of using (FIN) in the next SYN packet which forms a unusual behavior for the real user indeed and it generates a False positive (FP) value due to use of bloom filter data structure in case of congestion status. And it's inefficient in the case of Request Reply (RR) protocol that sends the CliACK in the next SYN request [21].

B. Flow Statistical Analysis

The analysis of statistical resulting in the packet tracing or packet counting and so on can be considered as packet sampling analysis from huge grown packets from any attacker source.Presents the approach that the proposed mechanism has the capability to detect SYN flooding attack accurately [9][24] .It has advantages as it uses threshold values to detect anomalies in flow rate which is determined on normal traffic statistics. [10]But it's not able to detect SYN flooding when the sampling rate is low. According to [10] [11] it has the advantages of low computational overhead because the proposed scheme doesn't hold the three way handshaking states but only analysis the SYN and ACK segments statistically but cannot overcome low rate SYN flooding attack.



C. AI detection scheme

It was adopted by many researches to design and implement intrusion detection system for denial of service attacks founded by [13]. According to [14] a system represented by two of the blocks to represent the packet classification and next one is the fuzzy system that determines the attacking possibility.

In the process the packet classification block classifies some of the packets from the networking incoming packets and checks if the fragmented offset value is zero or not. If it gets one then it's a SYN packet. [26] It checking process runs within a predetermined time interval and then it goes to the next block that is the fuzzy system which detects the SYNflooding. The proposed system [15] accuracy compared with the CUSUM for five attacks where in the high accuracy low false negative rate generates an earlier alarm using CUSUM algorithm. [16] It shows a good result using a fuzzy logic compared to CUSUM algorithm but it has disadvantages also. These are its difficult to model the traffic network before and after the attack due to linear and burst characteristics of packet flow moreover it depends on offset value in a TCP packet header which is a change due to network congestion and other states.

III. PROPOSED FLOOD DETECTION PROCESS

A. SYN Flood detection process

For detection the traditional method is used in the proposed system but in a different way. It can detect multiple BOTs at a time. Because it has a statistical method for reporting the incoming user accesses and determines its offset header. At the same time registers the user information in the database at the router. When it checks the incoming traffic at a determined time of Δt and divides the packets within it. For checking the behavior it checks the packets per second values. And then specify the attacking pattern if the attack is moving in a round robin process by comparing with the registered database a continuous message of attacked is generated in the server for every repeated SYNs. By the following mechanism the detecting process continues by studying the behavior of the traffic. The detecting process can be realized with a simplified form of pseudo flow chart in figure 2.

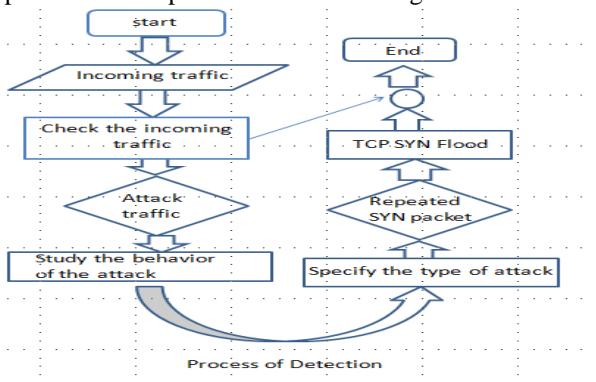


Fig. 2. Detection Process of TCP SYN Flood attack

In the systematic way if the system receives not more than 20 packets in the Δt time the process continues in the systematic way. Otherwise it will remain giving access to the SYNs and remove the entire database related with the accessed SYNs. In the proposed algorithm it's used the limited predetermined time interval of 20 seconds. In this experiment, it is chosen the second method to monitor the following is the algorithmic

process for detection. The algorithmic explanation helps in generation of the idea in this detection system.

B. Algorithmic Representation

Step 1: Accessing the network card of all access IP

Step 2: Explain the original IP of the access, record the time to obtain the access server.

Step 3: Find the historical IP access list with the currently accessed IP address, if, for each access to the IP and in the historical access list, the number of visits plus one. If the number of times is greater than 20, and the twentieth access time is earlier than the first if the access time is less than 20 seconds, it is determined that the current IP is the attacker IP, and the lostdiscard this IP and output the IP name.

//Determination of the behavior of packets and checking traffic.

Step 4: For each access of the IP which are not in the history access list, put this the new IP is added to the IP access. List and the number of accesses is 1, the first access time is the current system time.

//getting access for original IPs and resistive methods for spoofed one.

Step 5: Repeat the second step

Building a structure of Fake code instruction in brief
Struct {

Unsigned long ip; //attack ip

Intaccess_times; //attacks

Time access_first_time; // first attack time.

}iplist;

WHILE (1) {Ip = getacceessip();

IF ip IN iplist

THEN Set acess_times++;

 IF access_times<20 AND

 IF access_packet>20

THEN Print "You are attacked by IP"

ENF-IF };

Step 6: If condition fails sort database serially first unaccessed then accessed.

//making server ready for IP listing

Step 7: List the addresses of the attacker IP's.

Step 8: Repeat the process from step 1 followed by step 2.

//last step for getting the free port identity and IP

Addresses from database for accessing the defense process further.

In the implementing of the algorithm into a client system the resulting are much important for the studies in this flooding attack and in order to evaluate the detection sensitivity, there are two parameters being studied which are period to detect attacks initially and the period to take defensive measures. Based on the testing result of three attack patterns the above two parameters are summarized in the table 1. Obviously, the detection algorithm always firstly finds out the attackers, and then the victim is detected. However, fewer than three attack patterns, there is a very short interval between the periods to detect the attacker and victim. In the first attack pattern, we use a single attack source to attack the website and the traffic rate is 100 SYNs/s. The attack traffic rate in the third attack pattern is 100 times larger than that of the pattern one. The traffic rate of each individual attack source is also 10000 SYNs/s, but the attacker's number is virtually 100. The detection mechanism is not indifferent from the previous processes.

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It defines the attacking nature and in the graph there is shown the detecting time and defense time in versus positions. In the elaboration the defense time interval is nearby the detecting process. That is a good sign of the process as it runs simultaneously.

At the time of detection it starts automatic troubleshooting process that is described in the defense part. Recorded average experimental results for random packet attacks into a server and the recovery performance of the server itself are reported sequentially. For different values of generated packets and executing time is also noted from the system server operation by applying the detection and defense algorithms. An associated graphical representation is also reported as the attachment downward.

Table- I:Detection results under attack

SL No.	Attack Source	Total attack SYNs/S	Detecting time (min)	Defense time (min)
1	1	100	1.4	2.1
			3.5	2
2	10	10000	0.6	3
			1.4	2.1
3	100	50000	1.3	2.2
			1.1	2.1

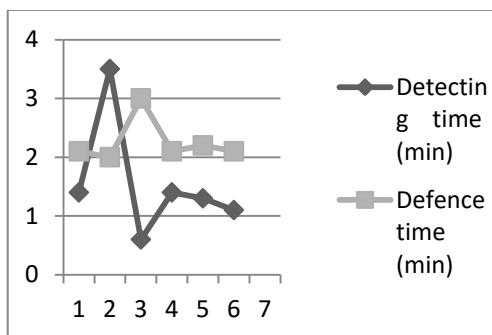


Fig. 3. Characteristic curve of simulation server under attack with proposed defense process.

IV. FLOODING RECOVERY PROCESS

A. Previous approaches

Following are the defense technique those are available in the SYN flooding defense based researches. The processes are followings according to [17].

B. SYN Cookies

In this technique no memory is reserved to store the initial request information on server. Instead a code is generated using the received initial request information and cryptographic techniques. This code is used as "sequence number" in the SYN/ACK packet and sent back [22]. When the respective ACK is received it extracts the initial request information and is used to set up the connection

C. SYNkill

Detects the attack due to SYN flood and then responds to lessen the effect of attack. SYNkill tool generates RST and ACK packets depending on the type of client request so that the resources of the server are not wasted.

D. Ingress Filtering

It is a source side defense system. It will only forward the packets that have the address of source (Prefix) same as the source network address (prefix). The exit or gateway routers are configured in such a way that they block all packets that are not the component of the input network address. It is more effective provided, it is implemented for all clients.

E. SYN Cache

In this method the resource allocation strategy is changed. Minimum initial information of the request is stored and then during the connection all the required information is stored.

F. SYNMON

Network processor is used as a processing unit to detect the SYN attack. An embedded system is designed to detect the attack using the CUSUM method.

G. A Router-based Novel Scheme

This method detects and mitigates the attack due to SYN flood. It is a router based mechanism that uses Bloom filter (counting) to keep track of number of SYN and FIN/RST packets. During the attack if the count of SYN is more than number of FINs, then the attack is detected. During mitigation every client's first request is dropped. The client retransmits the request and only such requests are forwarded to server. Thus the effect of attack is mitigated.

V. PROPOSED DEFENSE ALGORITHM

Following algorithm is generated as the prevention process of the SYN flood attacks. At the same time this algorithm can be used on large database systems. It is generally a module for the prevention of the DDoS attack can be used in general purpose server-client system. A cracking algorithm implemented to prevent SYN flooding attack by limiting the access of user. This helps to determine whether user is flooding attacker or legitimate user. When an attacker using genuine address, the proxy server uses the Deficit Round Robin algorithm to collect the address of the client request. If an attacker sends packets much faster than its fair share, the scheduling policy will drop its excess traffic. More Over, for each genuine IP address, the system will perform accounting on the number of packets that reach the firewall but are dropped by the scheduler; its IP address will be blacklisted. In some research papers the more efficient methodology is proposed to prevent SYN flooding attack by limiting the numbers of access to user or client [18]. The database is maintained between client and server which maintains the list of registered clients. So based on the database maintained the access is provided to registered users. In case of unregistered users the no of requests are checked and if threshold is not reached then access is granted. Also it depends on one more factor called — peak hours. During peak hours the request from the unregistered user is blocked temporarily.

Recovery process starts by executing its operation figured in the flowchart and algorithmic process of defense against the SYN flooding attacks that generated in the previous section. The defense activating process starts with the incoming traffic detection of a given time interval. And then it starts to operate with the incoming traffic. The database then accesses to operate its stored information about the attacks and drops the packets for stopping the attacks.

The algorithm and the process of simulation or programming code generation can be figured out by the following pseudo flow chart figure 4.

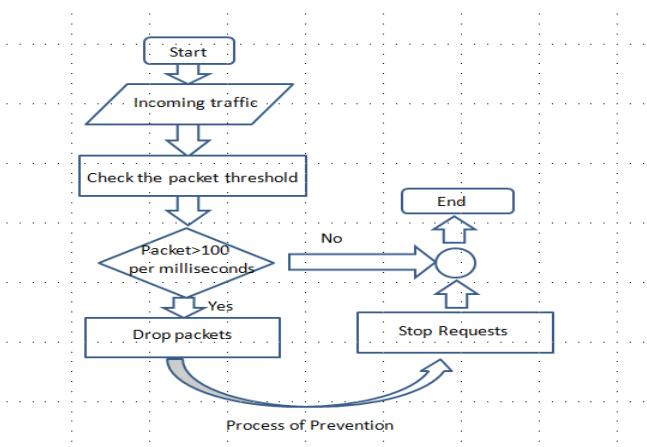


Fig. 4. Process of Defence and recovery of server.

A. Algorithm for defense mechanism

Following three terms in the system would be easier to find out the defending process of the algorithm influenced with [7].

- 1.) For preventing collect the physical address of a device.
- 2.) Match the network IP and the physical address of an IP.
- 3.) Count the number of Fake IPs of a physical address if ,

Step 1: Maintain the database for the list of users,X

Step 2: Analyze the User

Step 3: Get the useraddress of the incoming user.

If User=address of the incoming user Match it with the useraddress list in the database

Status="Registered"

Else

X.Login_Count++

Status="Unregistered"

End if

Next

// counting of available users into the Database.

Step 4: Response to the Request

If Status="Registered" then Process the Request and Send the Response

ELSE

If Status="Unregistered" then

Add name to the alert list, A

A.Name=User

A.Alert_count++

If A.Alert_count<Threshold_Value

If Server_peak_period=True

Add useraddress to Temp_Blocked List

Temp_Block=useraddress

End if

```

Else
Block the user permanently
P_Block=User
End if

//SYN-ACK response counting in fixed time
interval with respect to request from real user.

Step 5: List the detected useraddress
Step 6: Display the attacking useraddress
Step 7: Clear garbage database
Step 8: Get the useraddress of the new incoming user
Step 9: Repeat the Step 2 to Step 4
Step 10: Process the Request and the Response
  
```

Here the threshold value indicates the maximum request that a server may execute within a fixed time interval which a server provides as the peak period and if the packet exceeds the limit as directed in the condition of the flowchart as 100 packets per milliseconds then the server will catch flooding and this prevention process will operate its corresponding operation. In this algorithm the system is able to stop the flooding process by removing non registered addresses. By listing the addresses the system determines the amount of flooding. And then it starts to remove the garbage address information so no other flooding information can block space in the database.

B. SYN Flooding attack Detection and Defense Environment

In the early days these deadly attacks require environments like LINUX system for attacking. But now it can be implemented in various platforms. Several object oriented programming languages like JAVA, C++, PHP, C# are able to run this type of programs or modules for attacking. Creating BOTNETs from different PCs can also be done at the same time. Ubuntu 18.04 the updated version used for implementing the attacking client and affected server. As the BOTNETs generate random packets from different devices and the different devices generate spoofed addresses in IPv4. Separating the spoofed addresses the filter has to come forward to the server. At that time the filter has to check those spoofed address and find the real users from the requests. As long as the BOTNET requests become larger filter detects those ones and stop the requests immediately from those detected MAC addresses. For the Ubuntu 18.04 version there are some of basic commands that is followed for the simulation of the DDOS process. The GNU compiler and GCC runs the program attaching a thread header file as #include<pthread> and then the establishment of server client system access with C programming in LUNIX C. It accesses the SYN flooding from the client section. Multiple attacks can be generated in this process from different PCs. And server is able to defend all kinds of attacks within it.

C. Calculation for traffic in server client system

Packet classification determines efficiency of real-time detection and memory need, but extraction and storage of flood feature plays a significant role in flooding attack detection[6]. In fact, TCP, UDP and ICMP protocols are based on Client-Server model, and most of hosts in the network often play the dual role of client and server. Accordingly[6] [16],



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after analyzing the relevance between clients and servers, it shows that a strong synchronization exists between client's receiving SYN/ACK packets (SYN/ACK_rev) and it's sending ACK packets (ACK_sendkeep).

For DNS flooding attack and Smurf attack, the system can find out the attacks by checking the mismatch between the request packets and response packets. For client's traffic behavior detection, let $n \{ \Delta_n, n=1,2,3, \dots \}$ be the number of SYN/ACK_rev minus that of the corresponding ACK_sends collected within one detection period. And server's traffic behavior detection, let $\{ \Omega_n, n=1, 2, 3 \dots \}$ be the number of SYN_revs minus that of the corresponding SYN/ACK_sends and RST_sends collected within one detection period.

To alleviate the dependencies on the observation period, Δn is normalized by the average number F_1 of SYN/ACK_revs, and Ω_n is normalized by the average number F_2 of the sum of ACK_revs and RST_sends. F_1 and F_2 can be estimated in real time and updated periodically by the following recursive function.

$$F_2(n) = \mu F_2(n-1) + (1-\mu) ACK_rev(RST_send)(n) \dots \dots \dots \quad (2)$$

Where, η and μ are two coefficients strictly lying between 0 and 1. According to the highest weight of our past observations, μ and η are set to 0.7.....Define $\{X_n = \Delta n / F_1, n=1,2,3 \dots\}$ and $\{Y_n = \Omega_n / F_2, n=1,2,3 \dots\}$. So $\Sigma(X_n) = c(n-1)$ and $\Sigma(Y_n) = d(n-1)$. Moreover, $\{X_n\}$ $\{Y_n\}$ are no longer dependent on the network size and time-of-day. However, the CUSUM algorithm requires a negative drift before a change and positive drift after the change. We define $X_n = X_{n-\alpha}$ and $Y_n = Y_{n-\beta}$, where $\alpha > c$ and $\beta > d$. By this way, X_n and Y_n can fulfill the requirements of the CUSUM algorithm. During DDoS attack X_n and Y_n dramatically rise and become positive. But X_n and Y_n keep negative for normal traffic. Let

Where X_+ is equal to X if $X > 0$ and 0 otherwise. Y_n represents a stationary random process of client's traffic behavior, and Z_n represents that of server's traffic behavior. $d_n(\cdot)$ represents the decision at the end of every detection period, '0' for normal operation and '1' for attack (a change occurs). N_a and N_v respectively represent the detection threshold of the attacker and the victim $dna(n)=0$, if $Y_n \leq N_a$ Or $dna(n)=1$, if $Y_n > N_a$ $dvn(n)=0$, if $Z_n \leq N_v$ Or $dvn(n)=1$, if $Z_n > N_v$.

D. Simulation Outputs

For the associated process the algorithm needs to simulate this process properly. Following are some pictures where server defends the unauthorized requests. There is also a synchronization relationship between following three types of packets: server's receiving SYN packets (`SYN_rev`), its receiving ACK packets (`ACK_rev`) and sending RST packets (`RST_send`). Generally, most attackers choose many unreachable addresses as spoofed source addresses. Therefore, when a flooding attack starts, the strong correlation between the protocol handshake packets for TCP protocol will not occur.

```
bhagyo@bhagyo-HP-Pavilion-Notebook:~/Desktop/protik/source$ gcc -o client client.c -lpthread  
bhagyo@bhagyo-HP-Pavilion-Notebook:~/Desktop/protik/source$ sudo ./client 127.0.0.1  
[sudo] password for bhagyo: [REDACTED]
```

Fig. 5. For invoking the client section of the system by following commands.

In the system client IP address is fixed as we have generated. But in the BOTNET observations that this one can generate random spoofed addresses in IPv4. The server operation starts with two commands in the Ubuntu 18.04 terminal. Server detects the attack from the unauthorized user and after that takes action in the following way by using Prevention algorithm which is described earlier.

Fig. 6. Server under flooding attack

```
File Edit View Search Terminal Help  
239264 : you are attacked by lp 0.0.144.71  
239265 : you are attacked by lp 0.0.144.71  
239266 : you are attacked by lp 0.0.144.71  
239267 : you are attacked by lp 0.0.144.71  
239268 : you are attacked by lp 0.0.144.71  
239269 : you are attacked by lp 0.0.144.71  
239270 : you are attacked by lp 0.0.144.71  
239271 : you are attacked by lp 0.0.144.71  
239272 : you are attacked by lp 0.0.144.71  
239273 : you are attacked by lp 0.0.144.71  
239274 : you are attacked by lp 0.0.144.71  
239275 : you are attacked by lp 0.0.144.71  
239276 : you are attacked by lp 0.0.144.71  
239277 : you are attacked by lp 0.0.144.71  
239278 : you are attacked by lp 0.0.144.71  
239279 : you are attacked by lp 0.0.144.71  
239280 : you are attacked by lp 0.0.144.71  
239281 : you are attacked by lp 0.0.144.71  
239282 : you are attacked by lp 0.0.144.71  
239283 : you are attacked by lp 0.0.144.71  
239284 : you are attacked by lp 0.0.144.71  
239285 : you are attacked by lp 0.0.144.71  
239286 : you are attacked by lp 0.0.144.71
```

Fig. 7. Server attacked by the attacker.

The server starts listing the attacks after stopping that spoofed requests. Before taking actionserver lists those addresses. After that it resolves the Problems of the server while attacker client attacks to the server. It's like a shield, that standing by the server and defending the attacking activity. After using prevention algorithm it generates the following action in the terminal.

```
File Edit View Search Terminal Help
list ip ls 2831208395
list ip ls 889192575
list ip ls 134744072
list ip ls 16777343
list ip ls 16820416
list ip ls 0
list ip ls 100706496
list ip ls 35787
list ip ls 184592576
list ip ls 2831208395
list ip ls 889192575
list ip ls 134744072
list ip ls 16777343
list ip ls 16820416
list ip ls 0
list ip ls 100706496
list ip ls 35787
list ip ls 184592576
list ip ls 89513507
list ip ls 2831204845
list ip ls 2831208395
list ip ls 889192575
```

Fig. 8. Listed IPs by server working on recovery process

```

240014 you are ok to access with ip 11.0.168.192
240015 you are ok to access with ip 11.0.168.192
240016 you are ok to access with ip 11.0.168.192
240017 you are ok to access with ip 11.0.168.192
240018 you are ok to access with ip 11.0.168.192
240019 you are ok to access with ip 11.0.168.192
240020 you are ok to access with ip 11.0.168.192
240021 you are ok to access with ip 11.0.168.192
240022 you are ok to access with ip 11.0.168.192
240023 you are ok to access with ip 11.0.168.192
240024 you are ok to access with ip 11.0.168.192
240025 you are ok to access with ip 11.0.168.192
240026 you are ok to access with ip 11.0.168.192
240027 you are ok to access with ip 11.0.168.192
240028 you are ok to access with ip 11.0.168.192
240029 you are ok to access with ip 11.0.168.192
240030 you are ok to access with ip 11.0.168.192
240031 you are ok to access with ip 11.0.168.192
240032 you are ok to access with ip 11.0.168.192

```

Fig. 9. Server recovered the system after stopping the flooding attack

VI. COMPARISON AND PERFORMANCE DESCRIPTION

In the explanation of this proposed technique the system operates with both detection and defense against the SYN flooding attack, so the comparison must be with both algorithms. Comparison of the detection methods is described in the table II below

Table- II:Algorithm Comparison (for detection)

Technique	Method	Advantage	Disadvantage
Detection on edge router[20]	Network based router method	Guarantees that each packet sent by client is valid as much as possible.	Guarantees that each packet sent by client is valid as much as possible.
Statistical scheme[10]	Statistical analysis of traffic	Low false positive and false negative rate.-Short detection time	Cannot overcome the low-rate SYN flooding attack and consuming resources leads to shut down the available resources.
Adaptive Distributed Mechanism [19]	Machine learning	Faster detection and more accurate.	When a service is under attack, all traffic is to be blocked.
Chi square approach [19]	Statistical agent based intrusion detection system	Statistically analyze amount and variation of packet issued by the sender.	Limits the performance of communication because of the overhead in sending packets.
Proposed Approach	Time and Packet generation based tracking system..	Checks the databases and the registered list in it. Effective in real time attacking process.	Massive flow of one million of flooding can't be handled as it's an under developing project.

The related performance of the algorithm has several advantages and disadvantages. According to the previous described methods the approaches have not entered their

implementation by adding the database of incoming requests. And they claimed to predict the possibility of incoming request. But this system works in runtime evaluation and detection mechanism is better enough to detect the attacks at instance. And it can be operated in accessing large amount of packets. The proposed system doesn't leave any blockage while tracking the server. It does cope with the real time attacks affecting in a system. Following table describes the comparisons among the previous algorithms with the proposed one.

Table- III:Algorithm Comparison (for defense)

Method name	Method used	Results
Routerbased Novel scheme[8]	Counting Bloom filter to keep track of SYN and FIN/RST. Persistence of clientproperty.	Keeps track of number of SYNs and FIN/RSTs
SYN Cache[9]	Initial minimum information is stored	Percentage of connections set up and time taken during the attack
SYNkill[10]	Classification of client requests	Delay in setting up of connections and number of connections.
Proposed approach	Extraction of request behavior	Tracks the Flooding and lists the addresses of the flood attacking participants.

Resulting processes of defense has quite a bit problems that these are not much efficient in real time defense processes as the review research paper claims. Proposed process claims the flooding problem mitigation at a real time defense process. It mitigates the BOT client attacks and recovers the problem and moreover it manages space in the database for the new incoming users who are not registered but they have not the attacking SYNs as declared previously in the defense technique. However, it doesn't work in the low packet generation speed, and in maximum, for one hundred thousand packets it works relatively slow. Response of the system with or without has been calculated in the following Table 3. In this table it's described the performance of the system is better than the previous one.

Table- IV:Comparison with different systems

Experiments	Server Without using any System	Server With FPGA based protection system[6]	Server With the proposed System
Packet Transfer rate(bytes)	0.49	2.25	0.10
Time taken for test(in seconds)	825	64.179	60.002

VII. CONCLUSION

In the modern age the performance of any work is important for understanding the effectiveness of it.



Design and Implement a Real-Time Detection and Defence Mechanism Against the SYN Flood Attack in Server Client System

This work has a better performance than the related works as explained in the comparison section. The server client system is the medium for the sample implementation of the work; it can be broadly used elsewhere. The system server configured in LINUX (GNU) platform, the terminal display object is impressive and the run-time evaluation of the reading (values) is much accurate than traditional interpreter or any console outputs. The SYN flooding attack on TCP has been described in numerous of other publications as presented in aforesaid briefly. Generated tools for SYN flooding are available in coded forms now-a-days. Mitigating this tool's destructive section is a challenge for now. Several widely deployed operating systems implement the flooding tools and mitigation techniques but can't defend or detect successfully. That this operational document discusses about defeating SYN flooding attacks without removing any of the data and make server lines free to connect with another user within a least time interval. In at least some cases, these operating systems do not enable these counter measures by default; however, the mechanisms for defeating SYN flooding are well deployed, and easily enabled by end-users. In this paper the mitigation process is easier to implement in any appliances like IoT, IoE etc. Though it is now a module to operate the small system but the algorithm is effective for the several sectors. As the SYN flooding module is available in any programming platform so the system must be prepared to defend much deadliest attacks from any network or any confrontational flooding module is under the surveillance of future work of this paper. Detection and defense module in C programming is an informal implementation of the process as it can be implemented in any OOP, would be much effective. It should be taken care of.

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AUTHORS PROFILE



Sumonto Sarker has completed his BSc from Hajee Mohammad Danesh Science and Technology University and Masters from University of Dhaka. He attend a meritorious career in both while student and working. Now he is working as Assistant Professor in Hajee Mohammad Danesh Science and Technology University. He has developed several network related works, and has most interest in network security, Different type routing protocol about MANET, VANET and Under Water Network. He has personal skills in Networking, programming languages in C, C++, JAVA, Network Simulator (NS2.3 & NS 3), OMNET, MATLAB, database management system is one of his skills.
E-mail :sumonto@hstu.ac.bd





Kritimoy Bosuis is a student and researcher from Bangladesh, studied in Electronics and Communication Engineering from Hajee Mohammad Danesh Science and Technology University. He had a successful student life in early days and has personal skills in algorithm generation, database management and big data processing. He developed his programming skills in C,

C++, JAVA and Python. He has better skills in database management also. He worked for Database and framework related web apps. He has also game developing skills. Network security is one of his top listed interests. He also attended in different domestic programming contests. He has online problem solving records also.

E-mail: bosukritimov@gmail.com



Firdous Bin Ismail had a meritorious student life Who has studied in Electronics and Communication Engineering from Hajee Mohammad Danesh Science and Technology University? He has a personal experience of attending domestic programming contests. He has skills in algorithm generation, database management and programming skills in C,C++ and JAVA. He is mostly interested in Network security. He

has also participated in online programming contests and has a better skills in programming basics and database

. E-mail :bishalhstu@gmail.com



Md. Mahabub Hossain received B.Sc. and M.Sc. degrees from the Department of Applied Physics and Electronic Engineering, The University of Rajshahi, Bangladesh, and Ph.D. degree from the School of Electronics Engineering of Kyungpook National University, Republic of Korea. He worked as a postdoctoral researcher in Graduate School of

Engineering, Tohoku University, Japan in 2017-2018. Currently, he is a faculty member in the Department of Electronics and Communication Engineering, Hajee Mohammad Danesh Science and Technology University, Bangladesh. His research activities have primarily related to micro/nano-structures for computing processing, electronic systems and sensor applications.

E-mail:mahabub@hstu.ac.bd



Md. Mehedi Islam received B.Sc. and M.Sc. degrees from the Department of Applied Physics and Electronic Engineering, The Islamic University, Kushtia, Bangladesh. He attend a meritorious career in both while student and working. Now he is working as Associate Professor in the Department of Electronics and Communication Engineering, Hajee Mohammad Danesh

Science and Technology University. His research activities have primarily related to Optical Fiber Communication, Network Security, electronic systems and sensor applications

E-mail : mehedi@hstu.ac.bd